Lab2: User Pragrams

TA Session

Prepare1: Open up Pintos in your IDE/Text Editor

Prepare2: Launch Pintos container



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Some announcements:

- ➤ Lab 2 Code will due two weeks later
- ➤ No grace day
- ➤ Start early, Start early, Start early
- ➤ You can complete Lab2 from a clean codebase



Contents

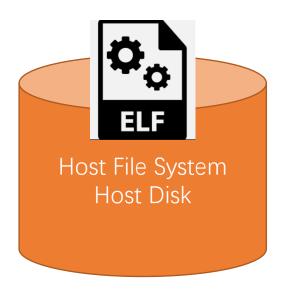
- Pintos Disk and File System
- System Call
- ➤ Interrupt Handling
- ➤ Lab2 tasks and suggestions

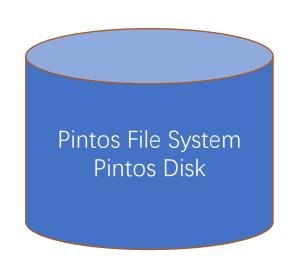


Where are the User Programs?

Source files are under /src/examples/ directory

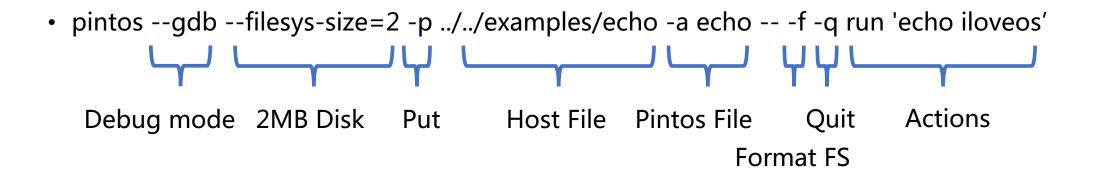
• Run `make` under /src/examples/ ELF files





Debug example:

Run `make && cd build` under /src/userprog/ directory



Details on Lab Document

Action: run 'echo iloveos'

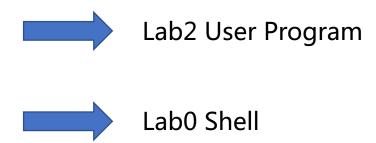
pintos_init():

```
/* Break command line into arguments and parse options. */
argv = read_command_line ();
argv = parse_options (argv);
```

```
printf ("Boot complete.\n");

if (*argv != NULL) {
    /* Run actions specified on kernel command line. */
    run_actions (argv);
} else {
    // TODO: no command line passed to kernel. Run interactively
}
```

```
(gdb) x/6s *argv
                 "-f"
0xc0007d3e:
0xc0007d41:
                 "-a"
0xc0007d44:
                 "extract"
0xc0007d4c:
                 "run"
                 "echo iloveos"
0xc0007d50:
0xc0007d5d:
(qdb) x/4s *arqv
0xc0007d44:
                 "extract"
0xc0007d4c:
                 "run"
0xc0007d50:
                 "echo iloveos"
0xc0007d5d:
```



Action: extract run 'echo iloveos'

run_actions():

```
/* An action. */
struct action
{
    char *name;
    int argc;
    void (*function) (char **argv);
};
```

```
/* Table of supported actions. */
  static const struct action actions[] =
      {"run", 2, run_task},
#ifdef FILESYS
      {"ls", 1, fsutil_ls},
      {"cat", 2, fsutil_cat},
      {"rm", 2, fsutil_rm},
      {"extract", 1, fsutil_extract},
      {"append", 2, fsutil_append},
#endif
      {NULL, 0, NULL},
    };
```

Action: extract run 'echo iloveos'

run_task():

```
/* Runs the task specified in ARGV[1]. */
static void
run_task (char **argv)
  const char *task = argv[1];
  printf ("Executing '%s':\n", task);
#ifdef USERPROG
  process_wait (process_execute (task));
#else
  run_test (task);
#endif
  printf ("Execution of '%s' complete.\n", task);
```

Q1: What is ARGV[0]?

Q2: What does run_test() do?

Action: run 'echo iloveos'

process_execute():

```
/* Create a new thread to execute FILE_NAME. */
tid = thread_create (file_name, PRI_DEFAULT, start_process, fn_copy);
if (tid == TID_ERROR)
   palloc_free_page (fn_copy);
return tid;
```

Main Thread:



Main Thread

run_task():

```
/* Runs the task specified in ARGV[1]. */
static void
run_task (char **argv)
  const char *task = argv[1];
  printf ("Executing '%s':\n", task);
#ifdef USERPROG
  process_wait (process_execute (task));
#else
  run_test (task);
#endif
  printf ("Execution of '%s' complete.\n", task);
```

Return Immediately!!



```
int
process_wait (tid_t child_tid UNUSED)
{
   return -1;
}
```

start_process:

```
/* Initialize interrupt frame and load executable. */
memset (&if_, 0, sizeof if_);
if_.gs = if_.fs = if_.es = if_.ds = if_.ss = SEL_UDSEG;
if_.cs = SEL_UCSEG;
if_.eflags = FLAG_IF | FLAG_MBS;
success = load (file_name, &if_.eip, &if_.esp);
```

- Load the ELF file from Disk into memory
- We are still in the kernel!!
- Initialize interrupt frame (eip, esp, segment registers, eflags)
- Start the user process by simulating a return from an interrupt

```
asm volatile ("movl %0, %%esp; jmp intr_exit" : : "g" (&if_) : "memory");
```

https://www.ibiblio.org/gferg/ldp/GCC-Inline-Assembly-HOWTO.html

load:

struct thread:

```
#ifdef USERPROG
    uint32_t *pagedir;
#endif
```

Create a new page directory:

```
/* Allocate and activate page directory. */
t->pagedir = pagedir_create ();
if (t->pagedir == NULL)
  goto done;
process_activate ();
                                    init_page_dir
                     NULL
t->pagedir =
                     Others
                                    Move t->pagedir into register CR3
```

PHYS BASE + 64MB

Kernel Mappings

```
user stack
                      grows downward
                       grows upward
            uninitialized data segment (BSS)
                initialized data segment
                      code segment
0x08048000
```

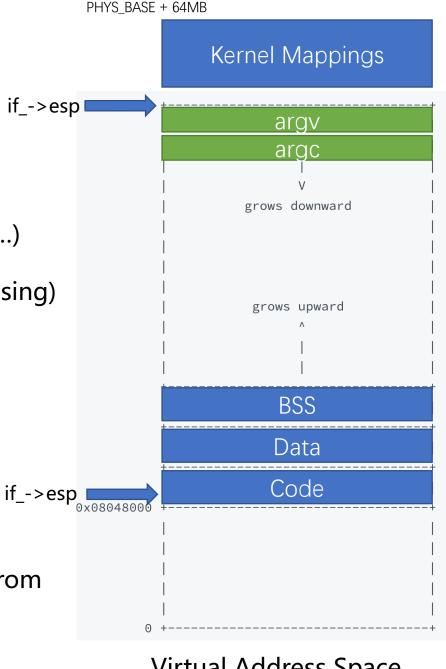
load:

- Read and verify ELF executable header
- Read ELF program header and load segments (code, data)
- Set up if_->stack (You will fix this in Exercise2: Argument Passing)
- Set up if_->eip with the entry point in executable header

/src/lib/user/entry.c:

```
void
_start (int argc, char *argv[])
 exit (main (argc, argv));
```

 After loading, start the user process by simulating a return from an interrupt with interrupt frame if_



Virtual Address Space

Wow, your process is running in User Space!

• But, we want system call support !!

/src/lib/user/syscall.h:

```
void
halt (void)
{
    syscall0 (SYS_HALT);
    NOT_REACHED ();
}
```

/src/lib/user/syscall.c:

```
/* Invokes syscall NUMBER, passing no arguments, and returns the
   return value as an `int'. */
#define syscall0(NUMBER)
          int retval;
          asm volatile
            ("pushl %[number]; int $0x30; addl $4, %%esp"
               : "=a" (retval)
               : [number] "i" (NUMBER)
               : "memory");
          retval;
```

Wow, your process is running in User Space!

But, we want system call support !!

80x86 Calling Convention!!

/src/lib/user/syscall.h:

/src/lib/user/syscall.c:

```
void
exit (int status)
  syscall1 (SYS_EXIT, status);
 NOT_REACHED ();
```

```
/* Invokes syscall NUMBER, passing argument ARGO, and returns the
   return value as an `int'. */
#define syscall1(NUMBER, ARG0)
          int retval:
          asm volatile
            ("pushl %[arg0]; pushl %[number]; int $0x30; addl $8, %esp"
               : "=a" (retval)
               : [number] "i" (NUMBER),
                 [arg0] "g" (ARG0)
               : "memory");
          retval;
        })
```

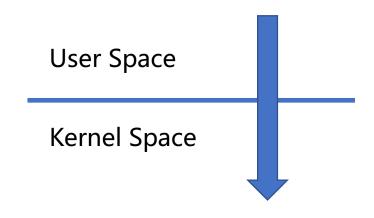
System Call Numbers:

/src/lib/syscall-nr.h

```
/* System call numbers. */
   /* Projects 2 and later. */
   SYS_HALT,
                             /* Halt the operating system. */
   SYS_EXIT,
                         /* Terminate this process. */
   SYS_EXEC,
                           /* Start another process. */
   SYS_WAIT,
                            /* Create a file. */
   SYS CREATE,
   SYS REMOVE,
                            /* Delete a file. */
                            /* Open a file. */
   SYS_OPEN,
   SYS FILESIZE,
                             /* Obtain a file's size. */
   SYS READ,
   SYS_WRITE,
                            /* Write to a file. */
   SYS SEEK,
                            /* Change position in a file. */
                             /* Report current position in a file. */
   SYS_TELL,
                             /* Close a file. */
   SYS_CLOSE,
   /* Project 3 and optionally project 4. */
   SYS_MMAP,
                    /* Map a file into memory. */
   SYS_MUNMAP,
                           /* Remove a memory mapping. */
   SYS CHDIR,
                             /* Change the current directory. */
   SYS MKDIR,
                            /* Create a directory. */
   SYS READDIR,
                             /* Reads a directory entry. */
                             /* Tests if a fd represents a directory. */
   SYS ISDIR,
   SYS_INUMBER
                             /* Returns the inode number for a fd. */
```

Now, all the magic is behind 'int 0x30'

```
/* Invokes syscall NUMBER, passing no arguments, and returns the
   return value as an `int'. */
#define syscall0(NUMBER)
        ({
          int retval;
          asm volatile
            ("pushl %[number]; (int $0x30; addl $4, %%esp"
               : "=a" (retval)
               : [number] "i" (NUMBER)
               : "memory");
          retval;
        })
```



Interrupt Handler

save the context of the interrupted

/src/threads/interrupt.c:

```
/src/user
```

```
void
intr_handler (struct intr_frame *frame)
```

```
void
syscall_init
{
  intr_regis
}
```

```
/* Invoke the interrupt's handler. */
handler = intr_handlers[frame->vec_no];
if (handler != NULL)
  handler (frame);
```

```
static void
syscall_hand
{
   printf ("s
   thread_exi
}
```

Implemen

```
/* Interrupt stack frame. */
struct intr_frame
   /* Pushed by intr_entry in intr-stubs.S.
      These are the interrupted task's saved registers. */
                               /* Saved EDI. */
   uint32_t edi;
   uint32_t esi;
                               /* Saved ESI. */
   uint32_t ebp;
                               /* Saved EBP. */
   uint32_t esp_dummy;
                               /* Not used. */
   uint32 t ebx;
                               /* Saved EBX. */
   uint32 t edx;
                               /* Saved EDX. */
   uint32 t ecx;
                               /* Saved ECX. */
   uint32_t eax;
                               /* Saved EAX. */
   uint16 t qs, :16;
                               /* Saved GS segment register. */
   uint16_t fs, :16;
                               /* Saved FS segment register. */
   uint16 t es, :16;
                               /* Saved ES segment register. */
                               /* Saved DS segment register. */
   uint16_t ds, :16;
   /* Pushed by intrNN_stub in intr-stubs.S. */
   uint32_t vec_no;
                               /* Interrupt vector number. */
   /* Sometimes pushed by the CPU,
      otherwise for consistency pushed as 0 by intrNN_stub.
      The CPU puts it just under `eip', but we move it here. */
                               /* Error code. */
   uint32_t error_code;
   /* Pushed by intrNN_stub in intr-stubs.S.
      This frame pointer eases interpretation of backtraces. */
                               /* Saved EBP (frame pointer). */
   void *frame_pointer;
   /* Pushed by the CPU.
      These are the interrupted task's saved registers. */
   void (*eip) (void);
                               /* Next instruction to execute. */
   uint16_t cs, :16;
                               /* Code segment for eip. */
   uint32_t eflags;
                               /* Saved CPU flags. */
   void *esp;
                               /* Saved stack pointer. */
   uint16_t ss, :16;
                               /* Data segment for esp. */
```

Contents

- Pintos Disk and File System
- System Call
- ➤ Interrupt Handling
- ➤ Lab2 tasks and suggestions



Some useful tips:

- Pintos exec == Unix fork + exec
- You can use malloc in kernel (#include "threads/malloc.h")
- Useful GDB command: loadusersymbols
- Reference to <u>xv6 implementation</u>
- multi-oom testcase will take some time, be patient

Step1: Argument Passing

/src/tests/main.c:

```
int
main (int argc UNUSED, char *argv[])
{
  test_name = argv[0];

  msg ("begin");
  random_init (0);
  test_main ();
  msg ("end");
  return 0;
}
```

Why stack?

- Set up the stack after loading
- Argument Passing details in Lab Doc
- Pass all the args-xxx tests

Step2: Halt System Call

- Argument Passing
- System Call Infrastructure

Step3: Some temporal workaround

- The exit system call (barely work is fine)
- The write system call for writing to fd 1, the system console
- change process_wait() to an infinite loop (one that waits forever)

Step4: Accessing User Memory

- User programs will pass arguments (char*, int, unsigned) into kernel
- These arguments are on the user stack (esp is saved in intr_frame)
- Ensure the address validity (in user page table)
- Avoid repeating code !!
- Two implementation suggestion (in Lab Doc)

Step4: Process Control System Call

- exit, exec, wait
- Design all at first, they may share some data structures

Step5: FS System Call

- No need to understand file system implementation
- Read the interfaces in /src/filesys/file.c, /src/filesys/filesys.c
- Pass all tests but rox-simple, rox-child, rox-multichild

Step7: Denying Writes to Executables

- Why?
- Close a file will re-enable writes
- Keep the executable file open during execution

Step8: Cheers !!!



Enjoy Your Pintos Journey ~~

Any Problem?